



ZeRØ: Zero-Overhead Resilient Operation Under Pointer Integrity Attacks

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COLUMBIA | ENGINEERING

The Fu Foundation School of Engineering and Applied Science

About Me



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Memory Safety is a serious problem!

Computing Sep 6

...

Apple says China's Uighur Muslims were targeted in the recent iPhone hacking campaign

The tech giant gave a rare statement that bristled at Google's analysis of the novel hacking operation.

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EDITOR'S PICK | 42,742 views | Nov 21, 2018, 07:00am

Exclusive: Saudi Dissidents Hit With Stealth iPhone Spyware Before Khashoggi's Murder

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The New York Times

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WhatsApp Rushes to Fix Security Flaw Exposed in Hacking of Lawyer's Phone

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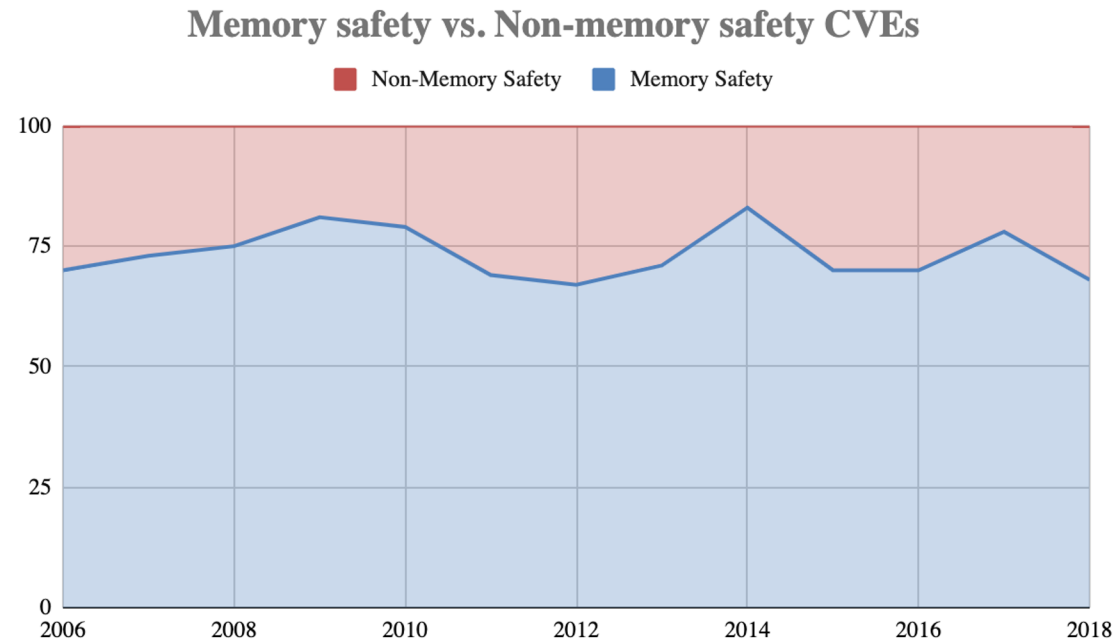
It's easy to make mistakes

It's easy to make mistakes



SEGFAULT!

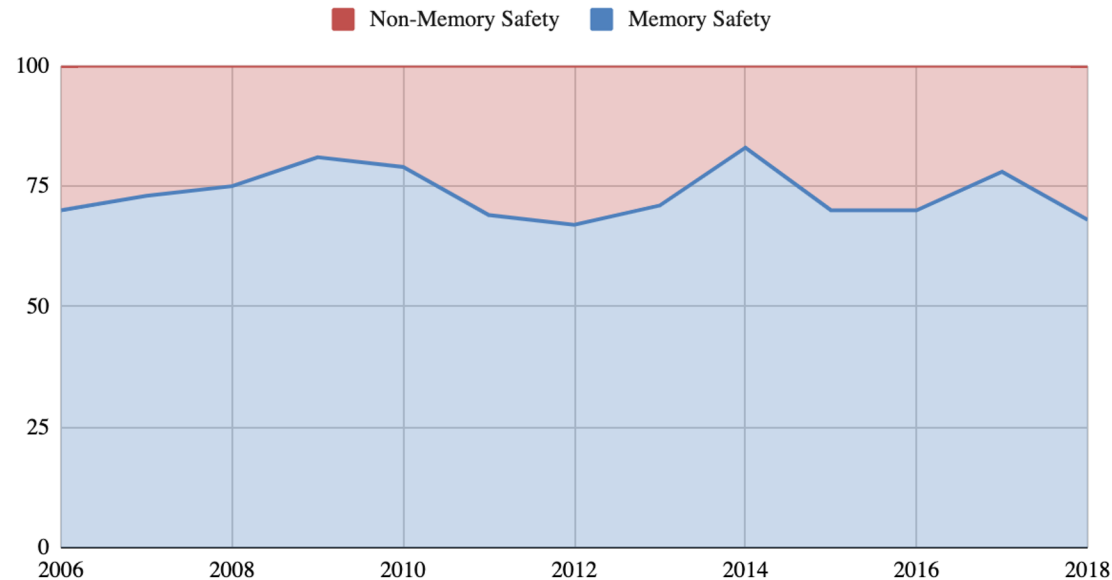
Prevalence of Memory Safety Vulns



Microsoft Product CVEs

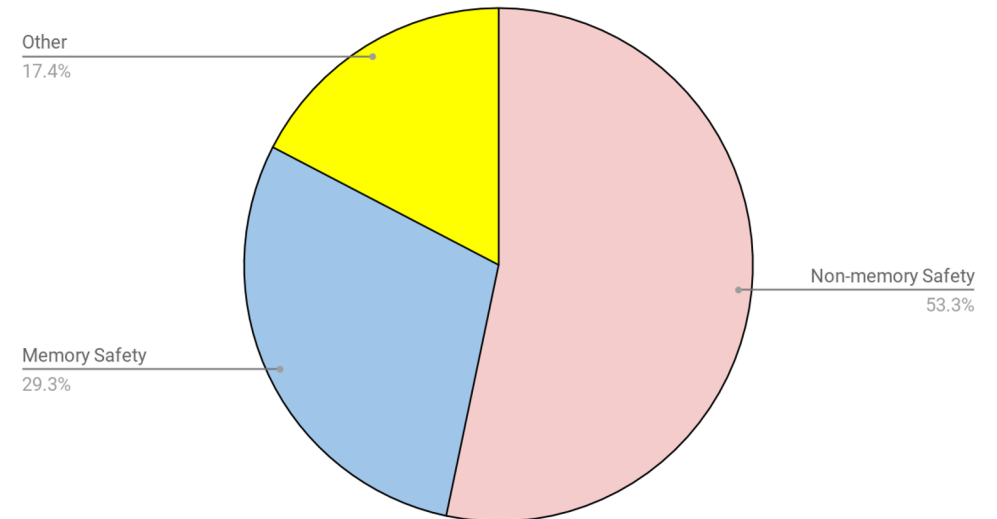
Prevalence of Memory Safety Vulns

Memory safety vs. Non-memory safety CVEs



Microsoft Product CVEs

OSS-Fuzz Bug Types



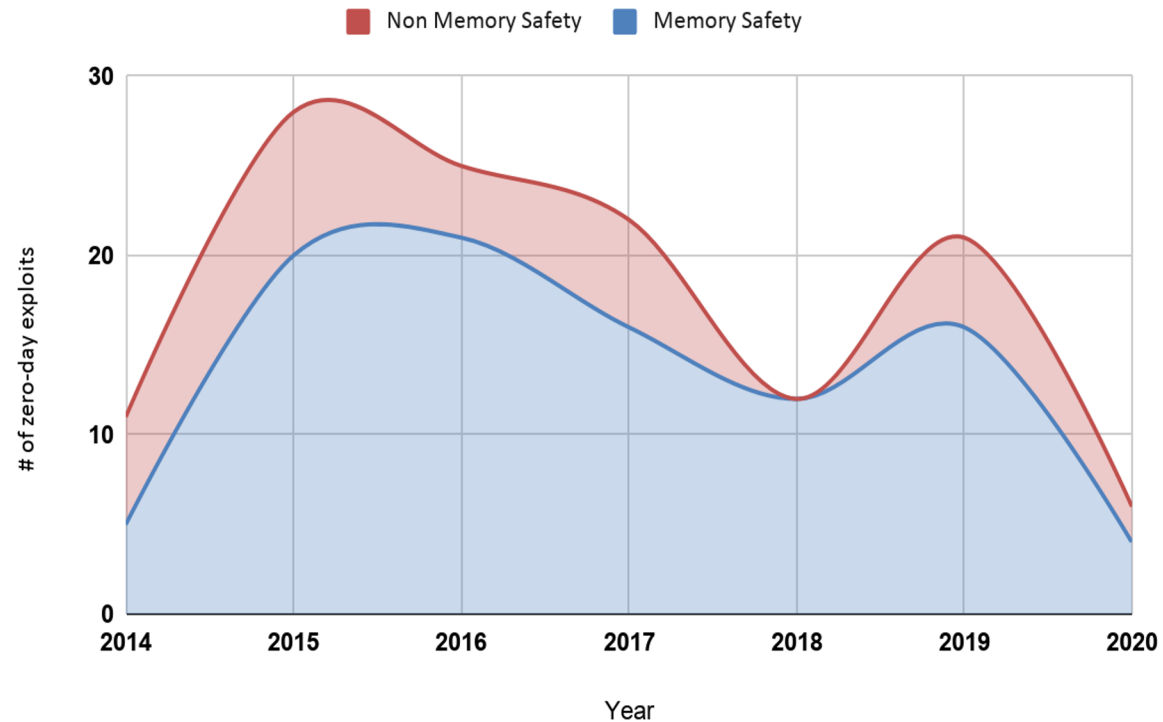
Google OSS-Fuzz bugs from 2016-2018.

ATTACKERS



MEMORY SAFETY

Attackers Prefer Memory Safety Vulns

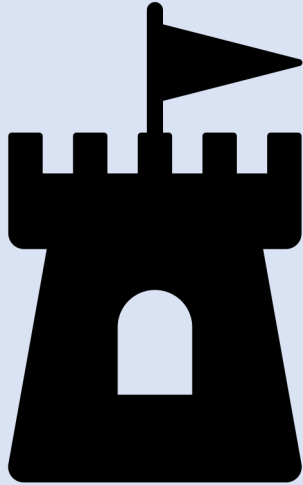


Zero-day “in the wild” exploits
from 2014-2020

How To Fix Memory (Un)Safety?

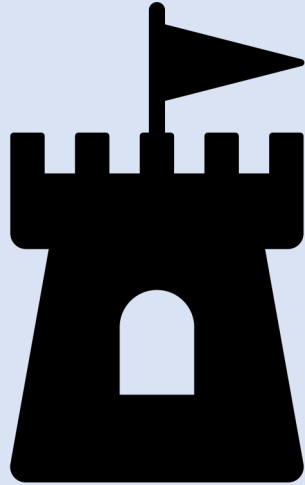
How To Fix Memory (Un)Safety?

**Memory Safe
Languages**



How To Fix Memory (Un)Safety?

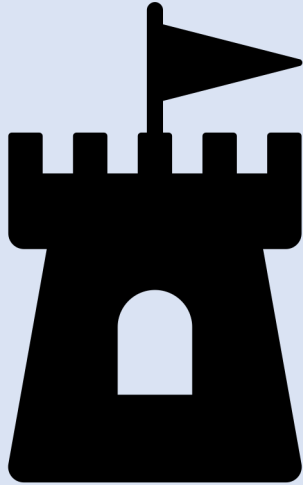
Memory Safe Languages



- Performance?
- Legacy Code?

How To Fix Memory (Un)Safety?

Memory Safe Languages



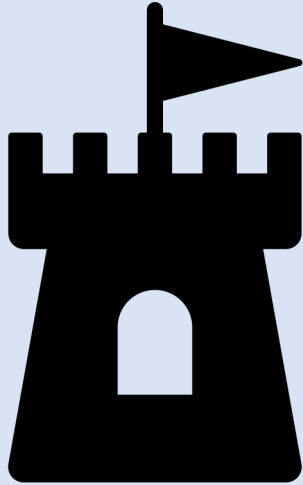
- Performance?
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Pre-deployment Testing



How To Fix Memory (Un)Safety?

Memory Safe Languages



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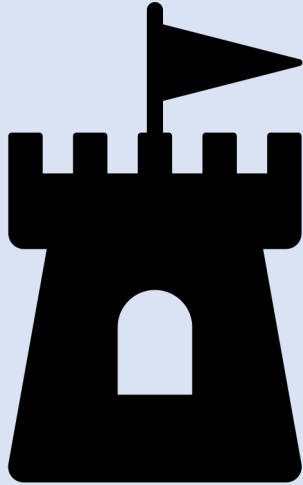
Pre-deployment Testing



- Time?
- Scalability?

How To Fix Memory (Un)Safety?

Memory Safe Languages



- Performance?
- Legacy Code?

Pre-deployment Testing



- Time?
- Scalability?

Post-deployment Mitigations

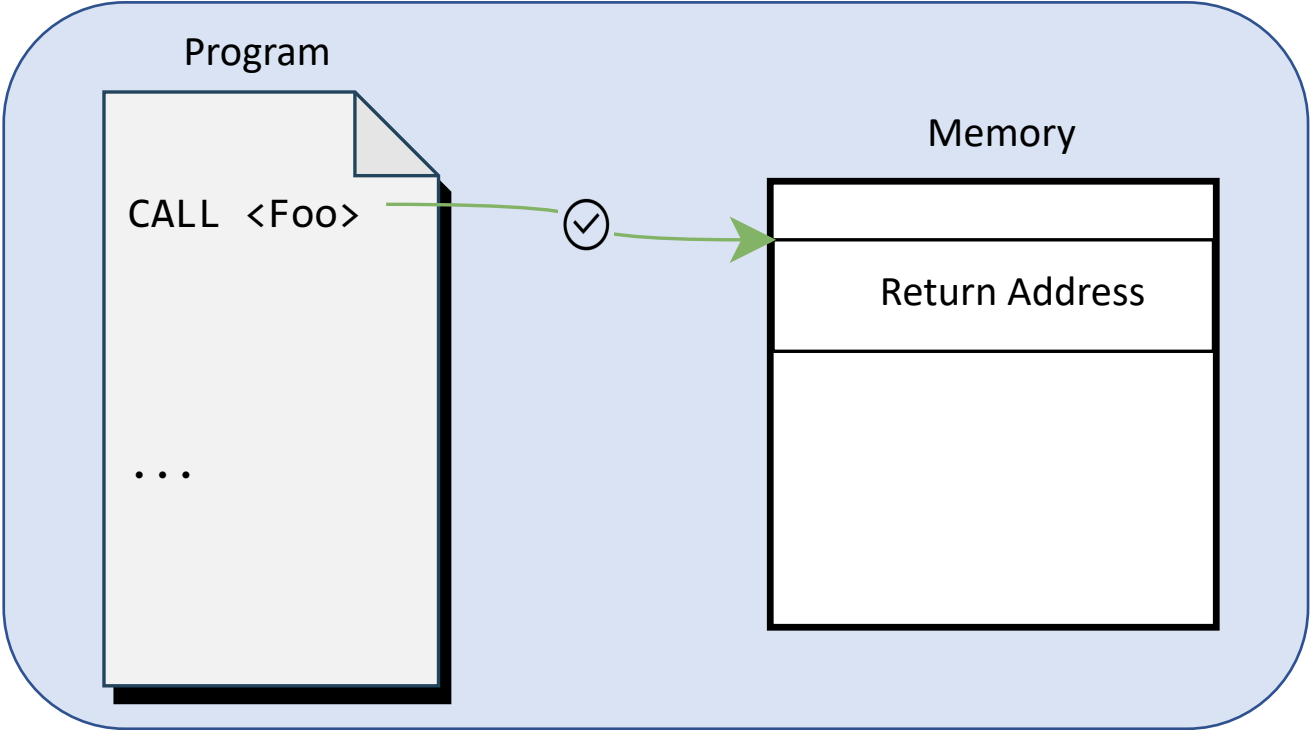




ZeRØ

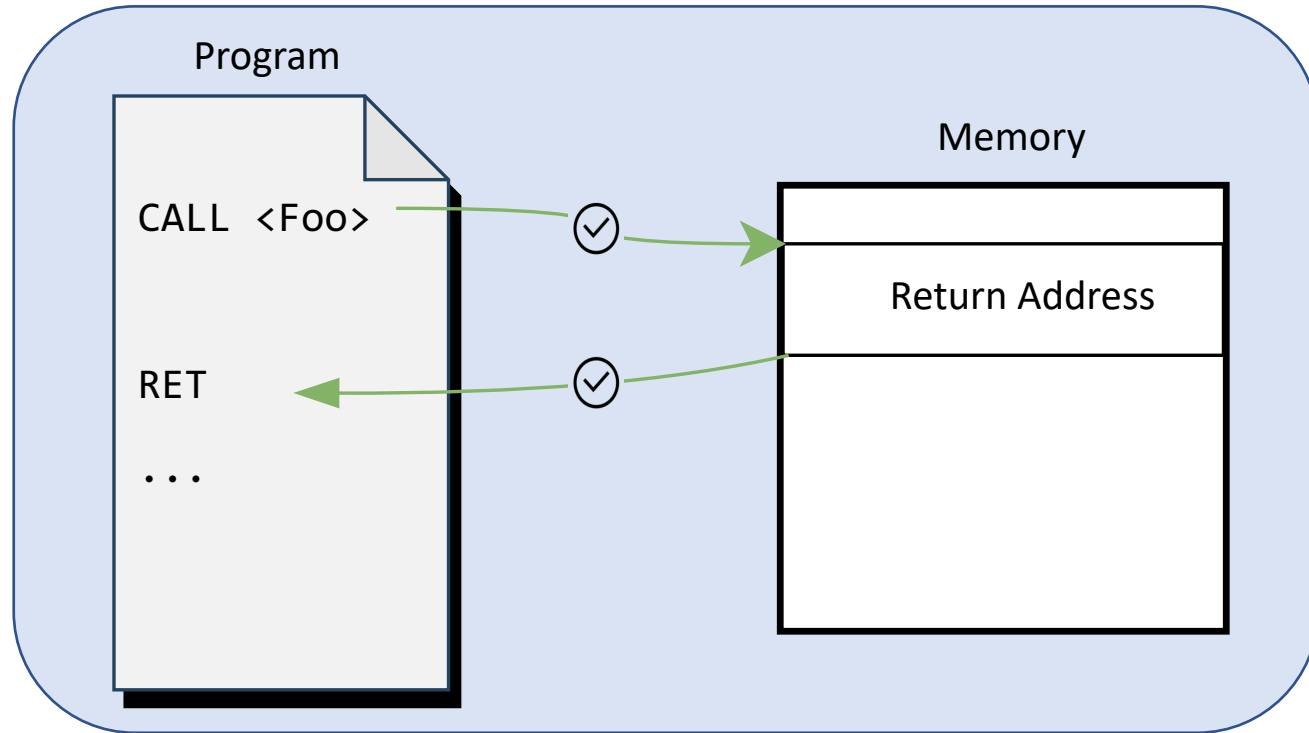
Overview

ZeRØ: Overview



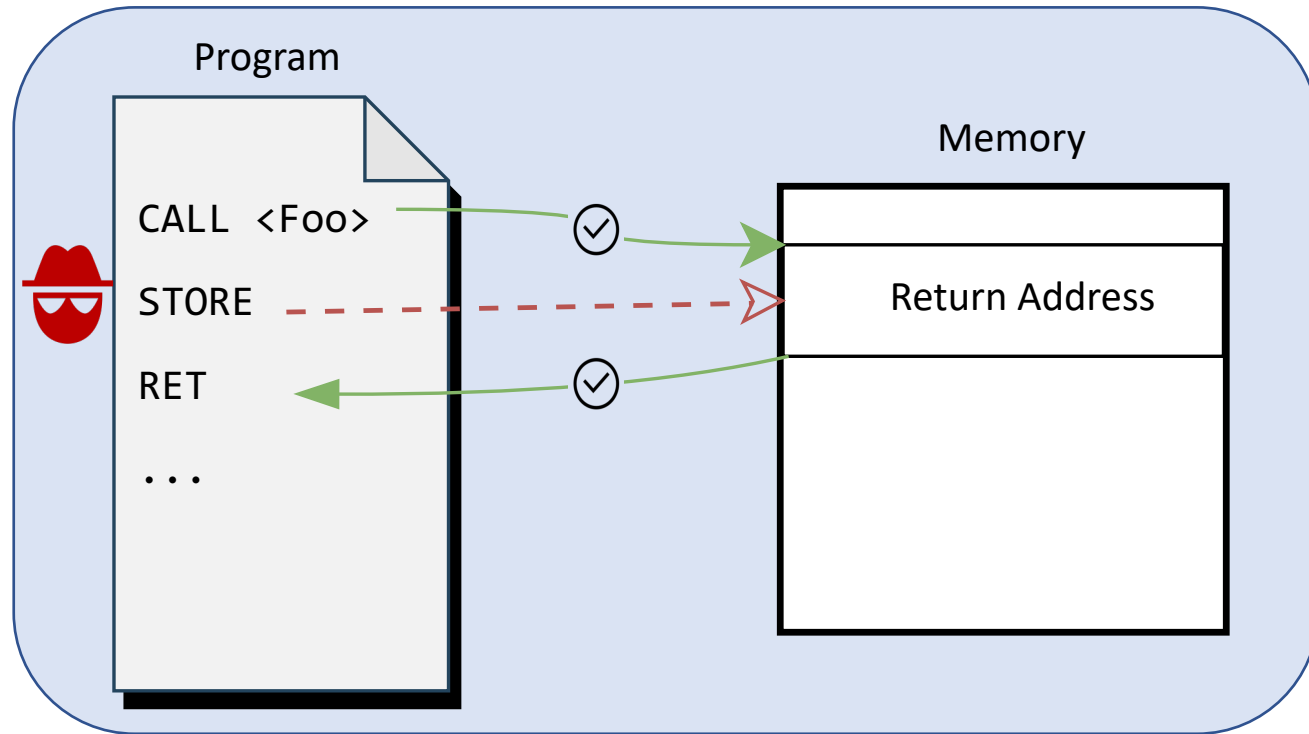
Return Address Protection

ZeRØ: Overview



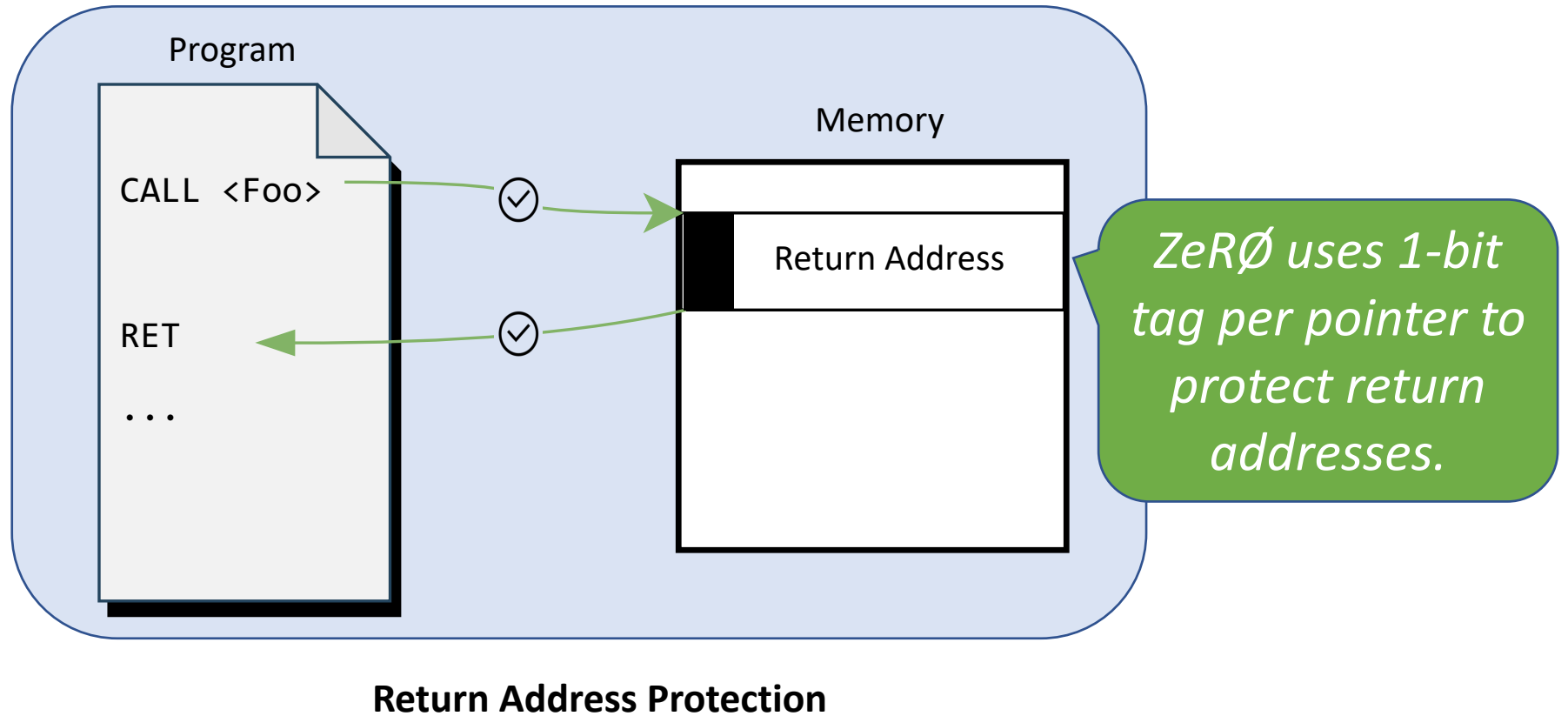
Return Address Protection

ZeRØ: Overview

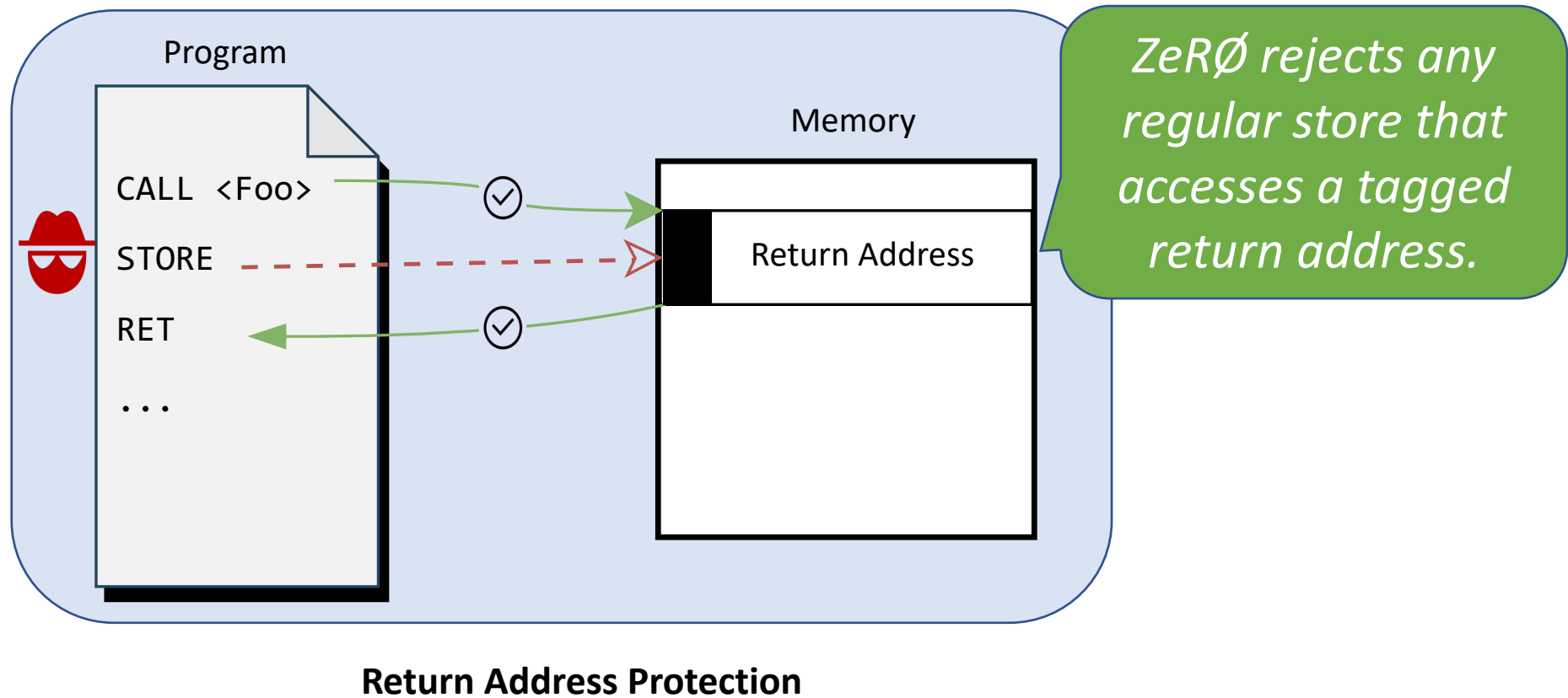


Return Address Protection

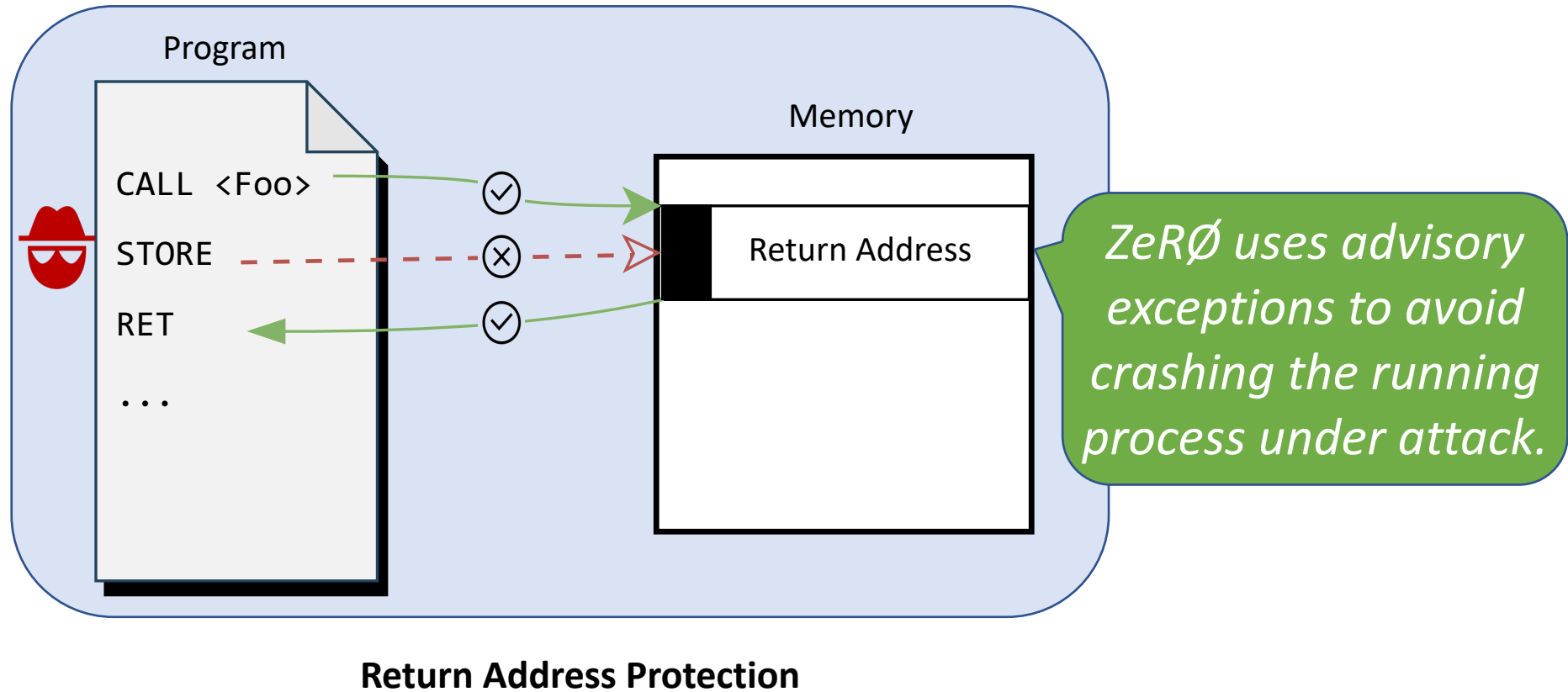
ZeRØ: Overview



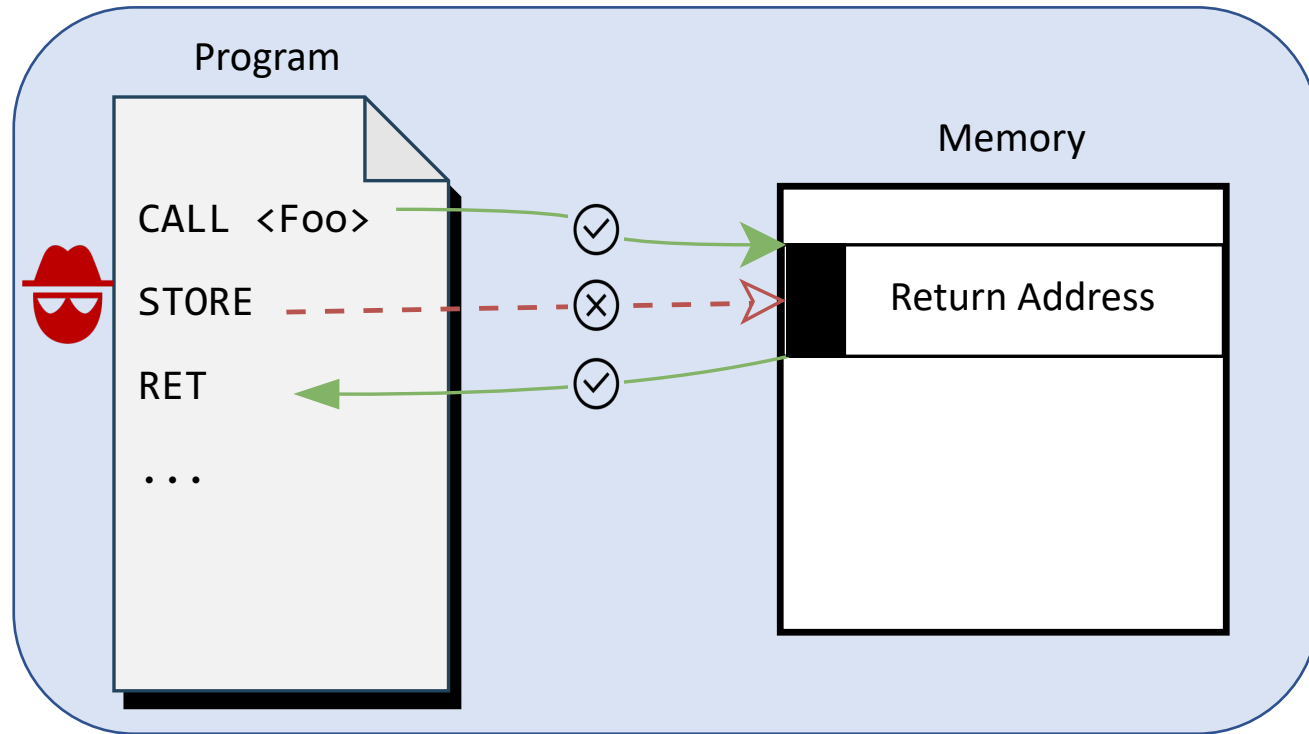
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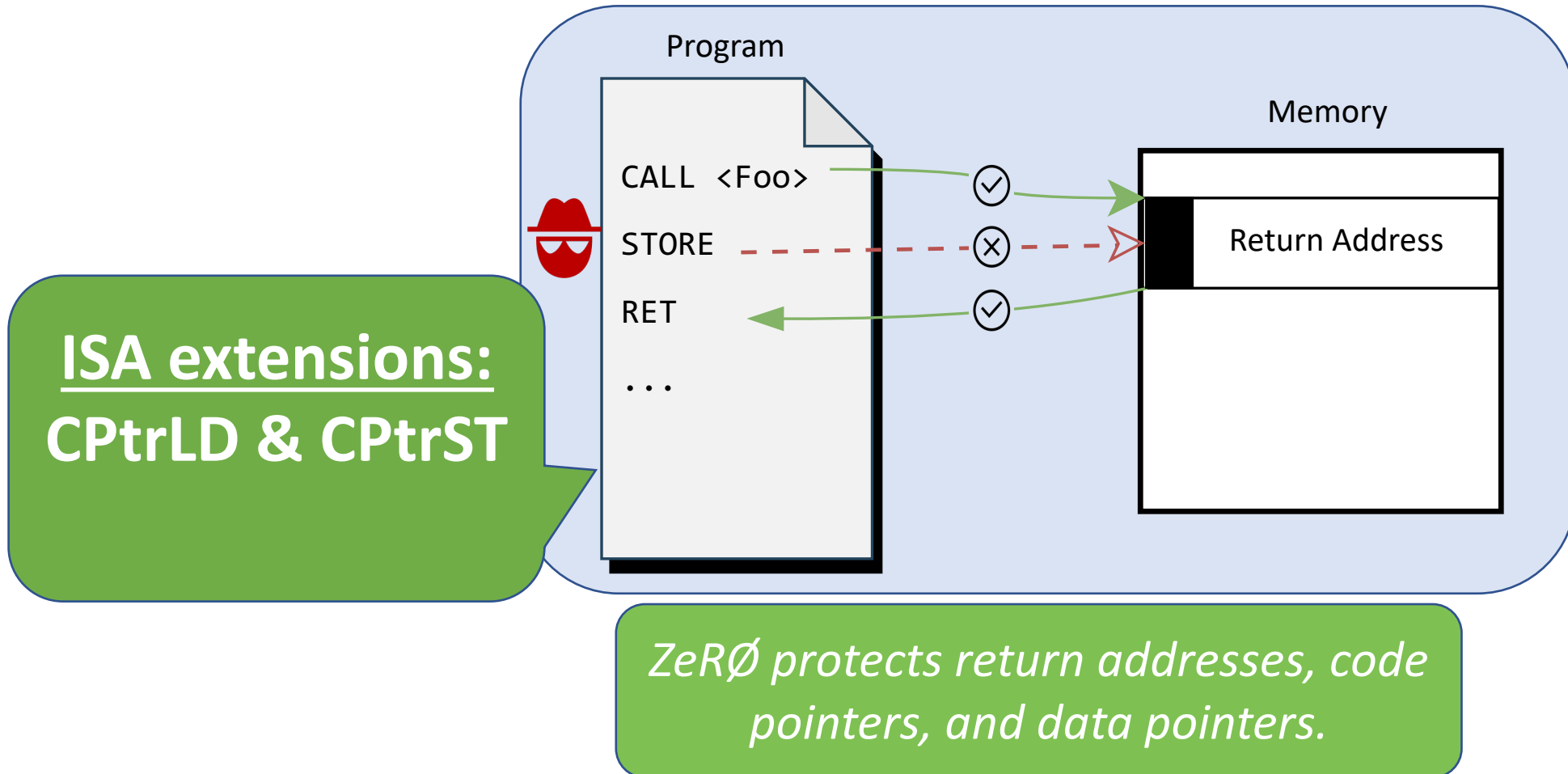


ZeRØ: Overview

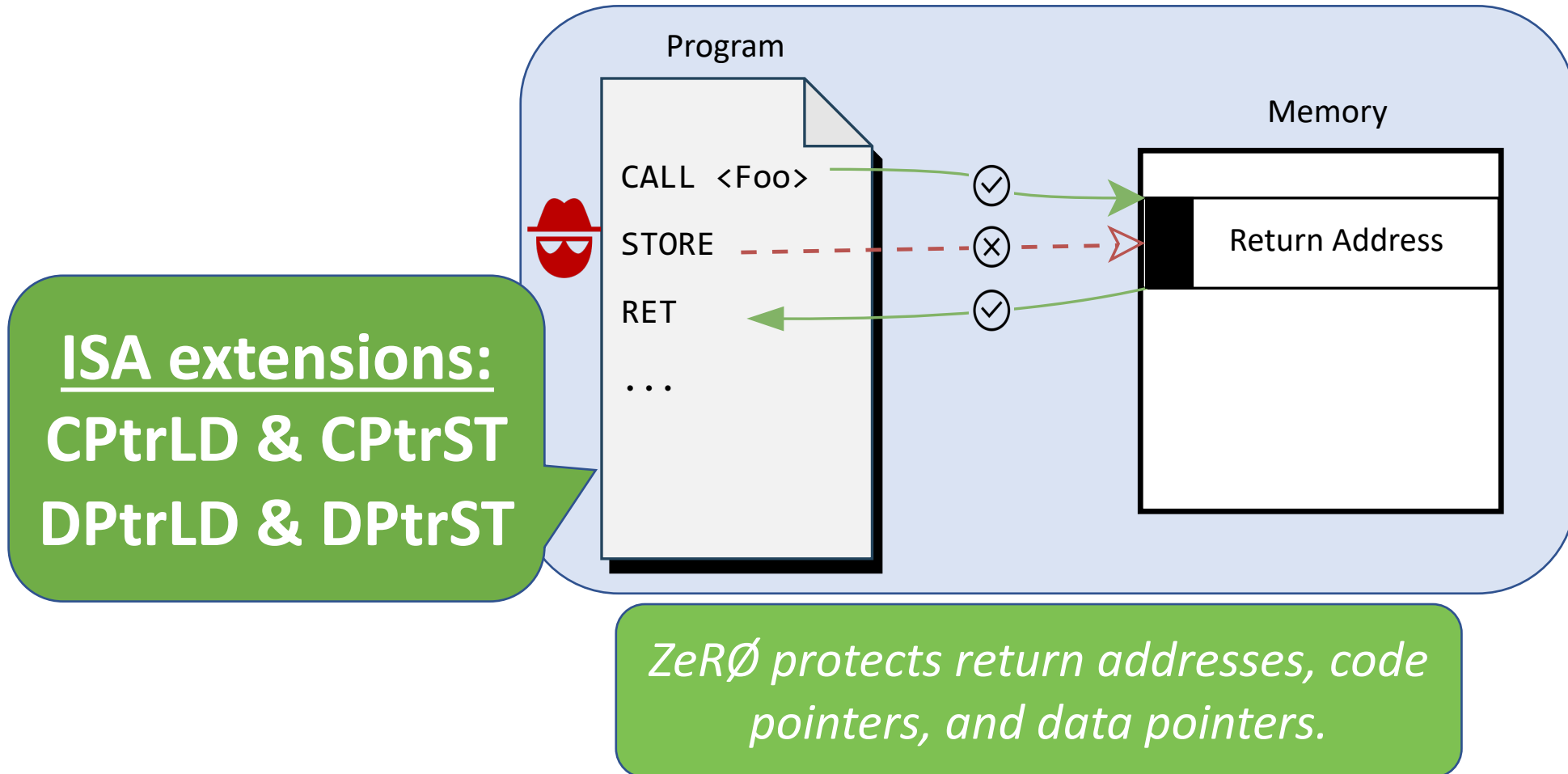


ZeRØ protects return addresses, code pointers, and data pointers.

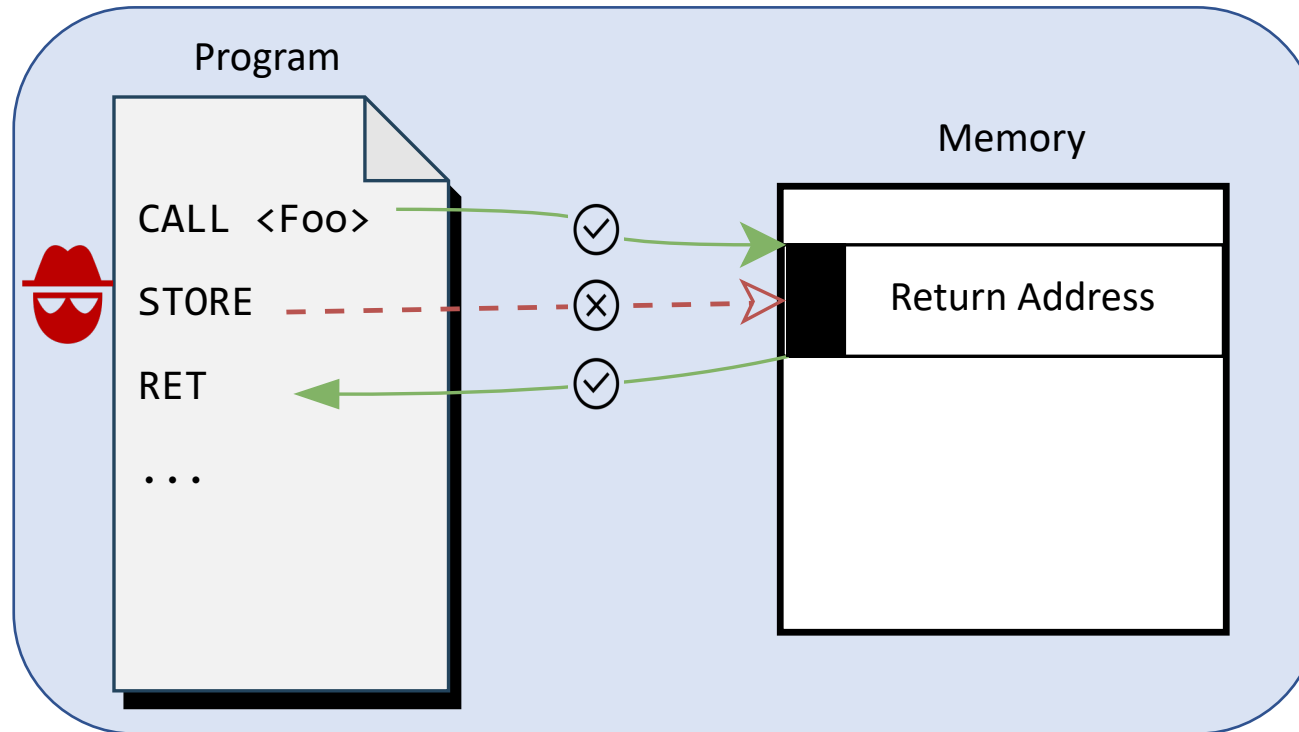
ZeRØ: Overview



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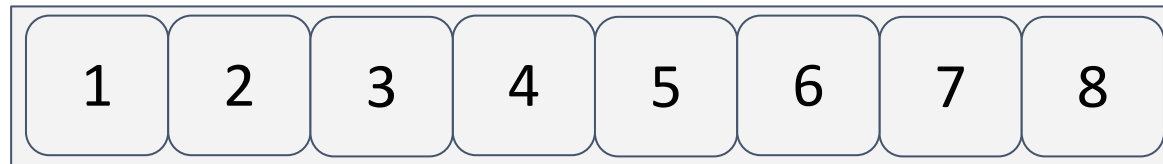


How can ZeRØ efficiently identify if a memory word is a return address, code pointer, data pointer, or regular data?

Cache Line Formats

ZeRØ: Cache Line Formats

Normal



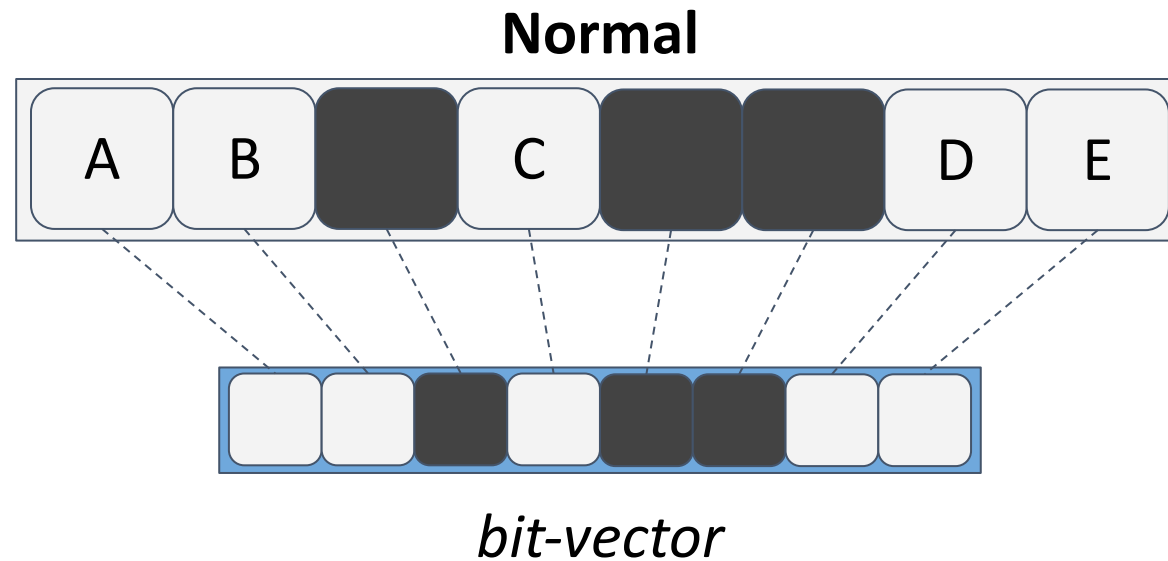
ZeRØ: Cache Line Formats

 Program
Pointers

Normal



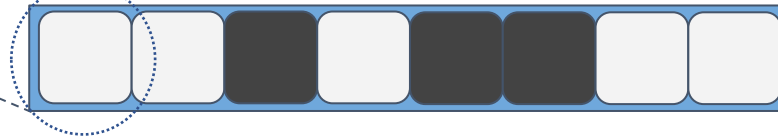
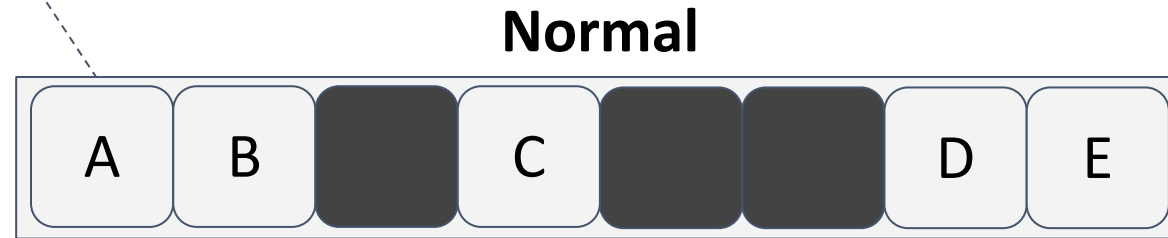
ZeRØ: Cache Line Formats



ZeRØ: Cache Line Formats

■ Program Pointers

Type	Bits
Return address	01

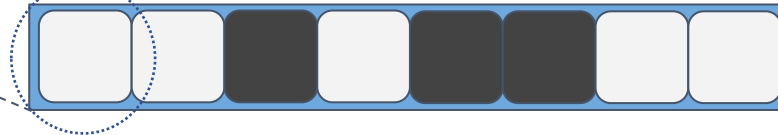
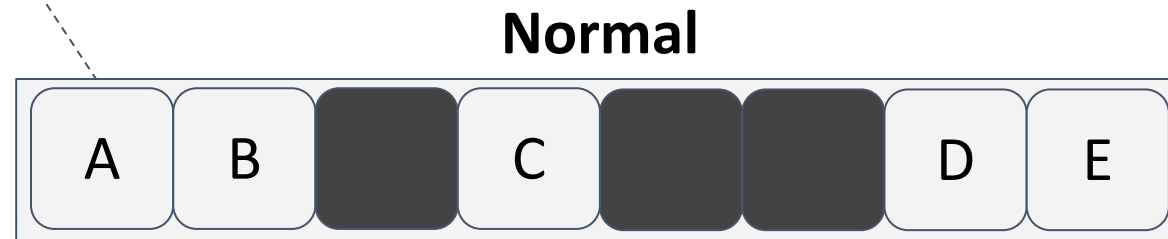


bit-vector

ZeRØ: Cache Line Formats

■ Program Pointers

Type	Bits
Return address	01
Function pointer	10

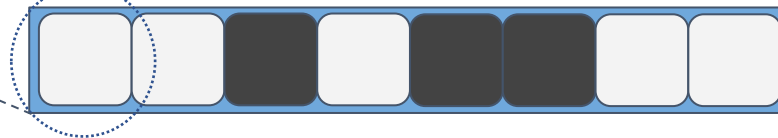
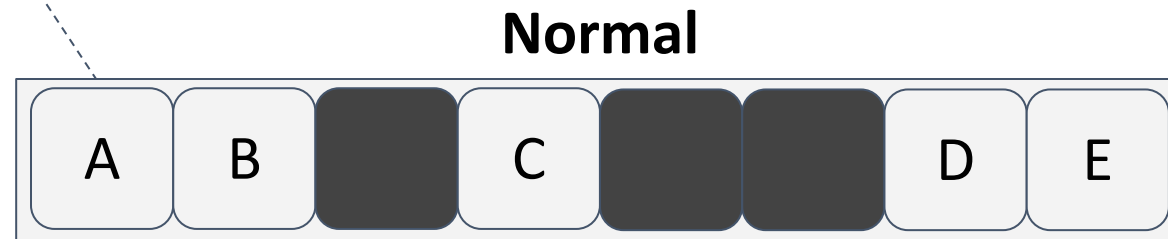


bit-vector

ZeRØ: Cache Line Formats

■ Program Pointers

Type	Bits
Return address	01
Function pointer	10
Data pointer	11

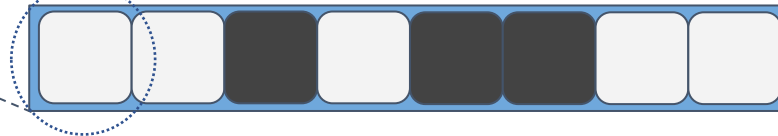
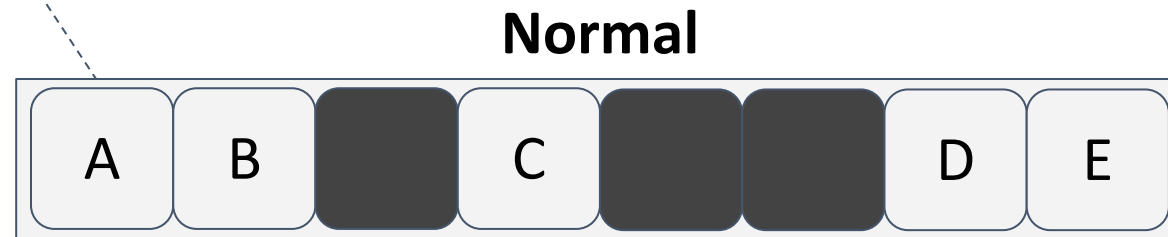


bit-vector

ZeRØ: Cache Line Formats

■ Program Pointers

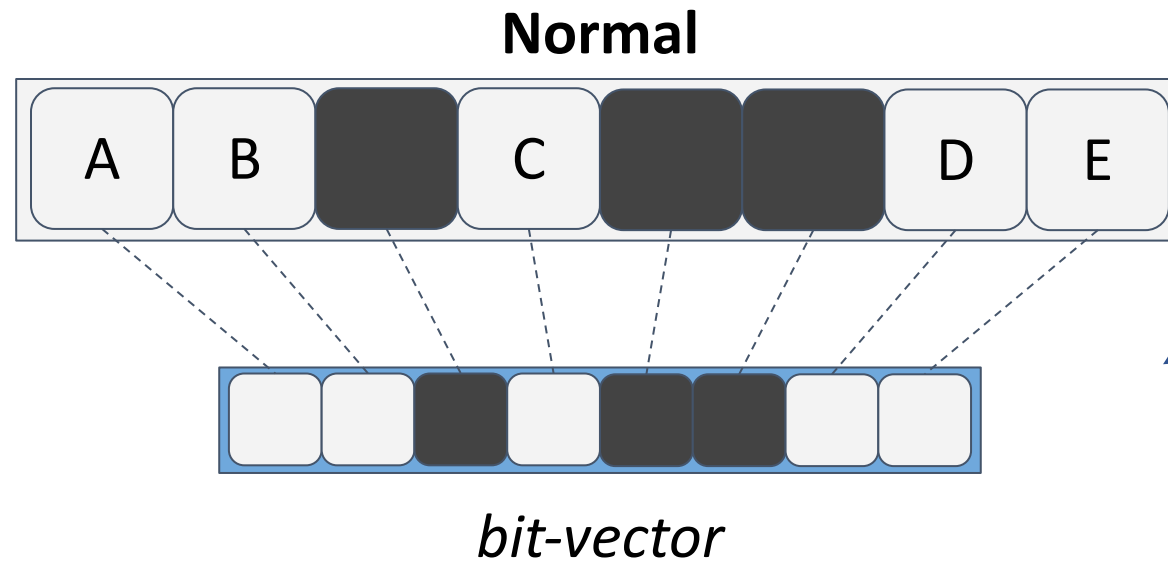
Type	Bits
Regular data	00
Return address	01
Function pointer	10
Data pointer	11



bit-vector

ZeRØ: Cache Line Formats

■ Program
Pointers



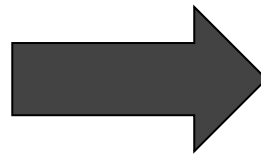
*3.125% area
overhead*

ZeRØ: Cache Line Formats

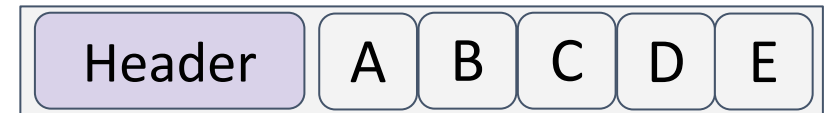
Our Metadata: Encoded within unused pointer bits.

■ Program
■ Pointers

Normal



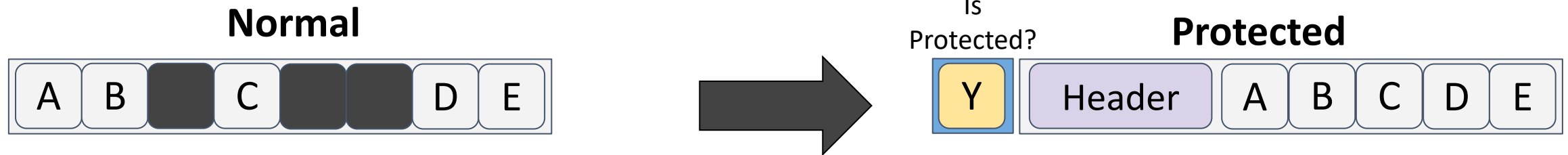
Protected



ZeRØ: Cache Line Formats

Our Metadata: Encoded within unused pointer bits.

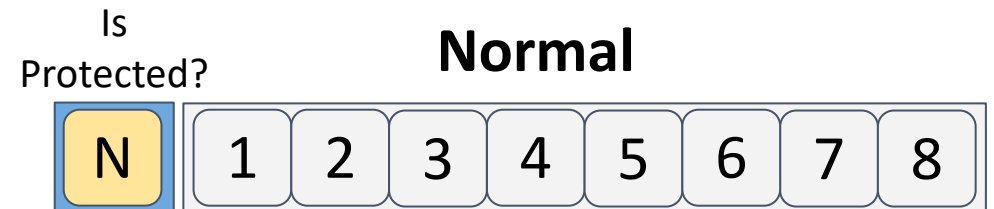
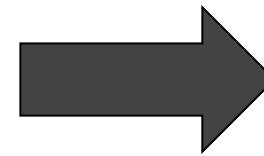
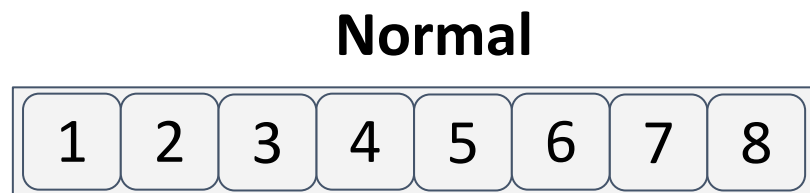
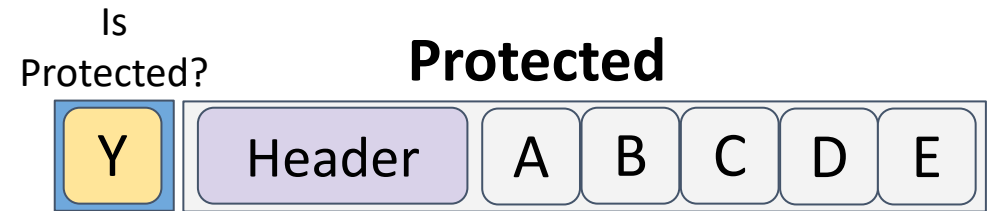
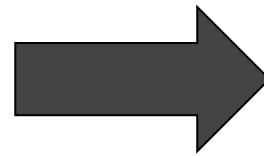
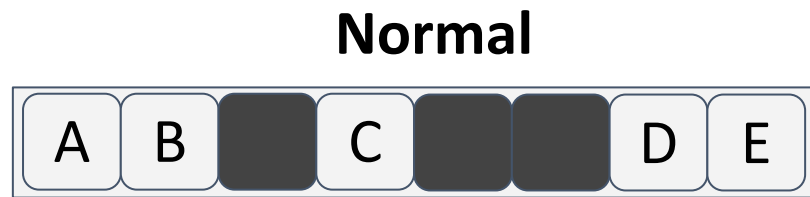
■ Program
■ Pointers



ZeRØ: Cache Line Formats

Our Metadata: Encoded within unused pointer bits.

■ Program
Pointers

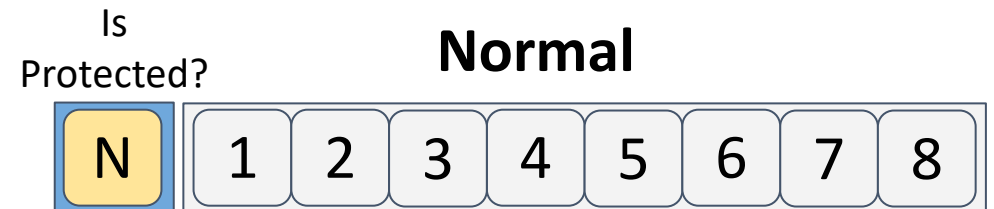
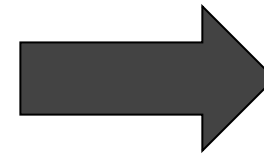
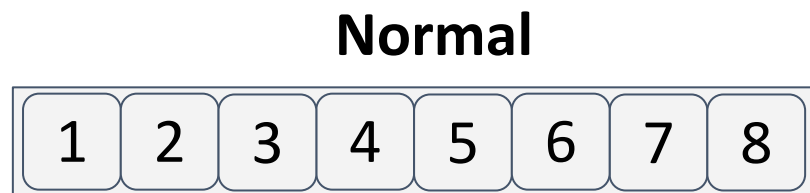
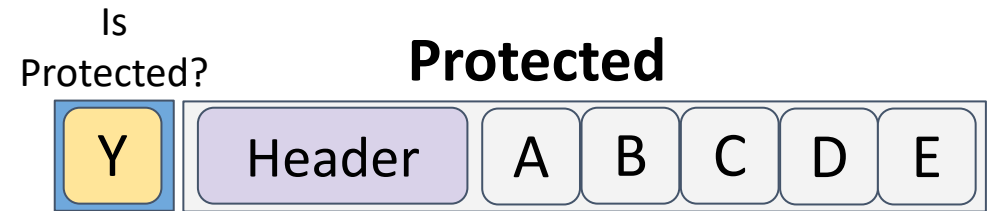
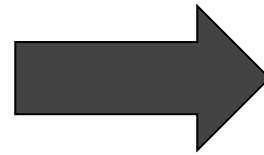
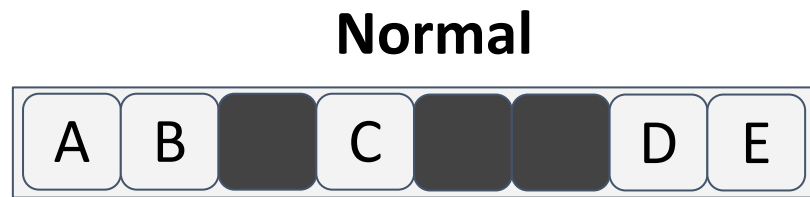


ZeRØ: Cache Line Formats

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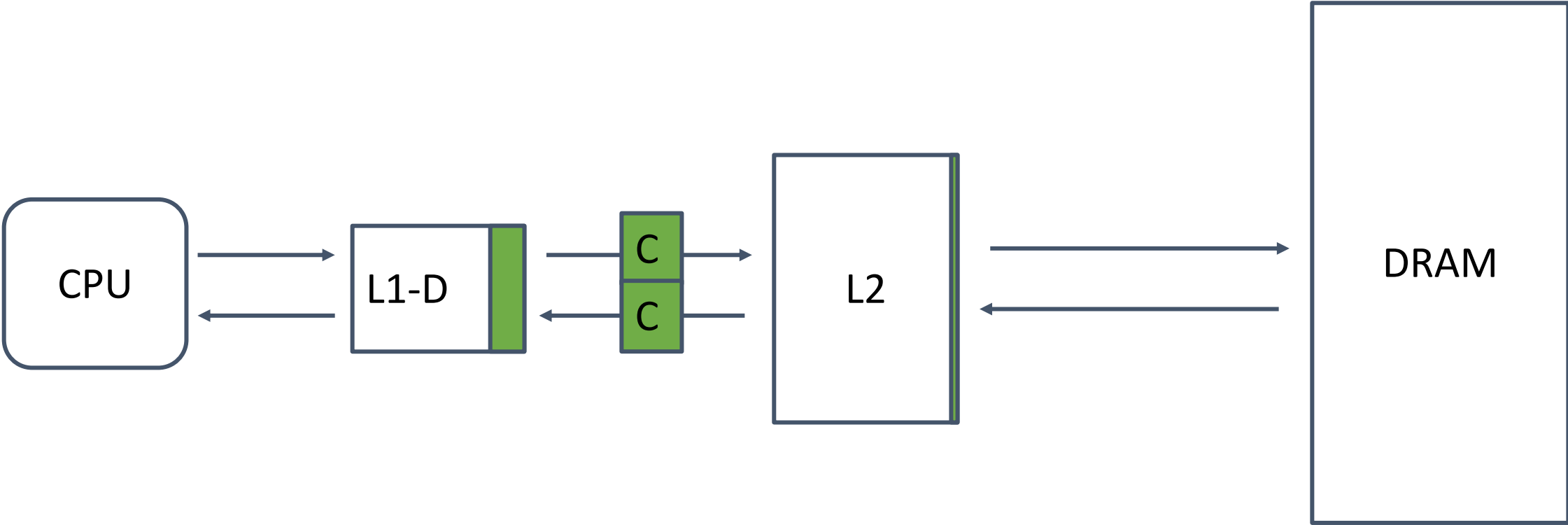
 Program
Pointers

*0.2% area
overhead*

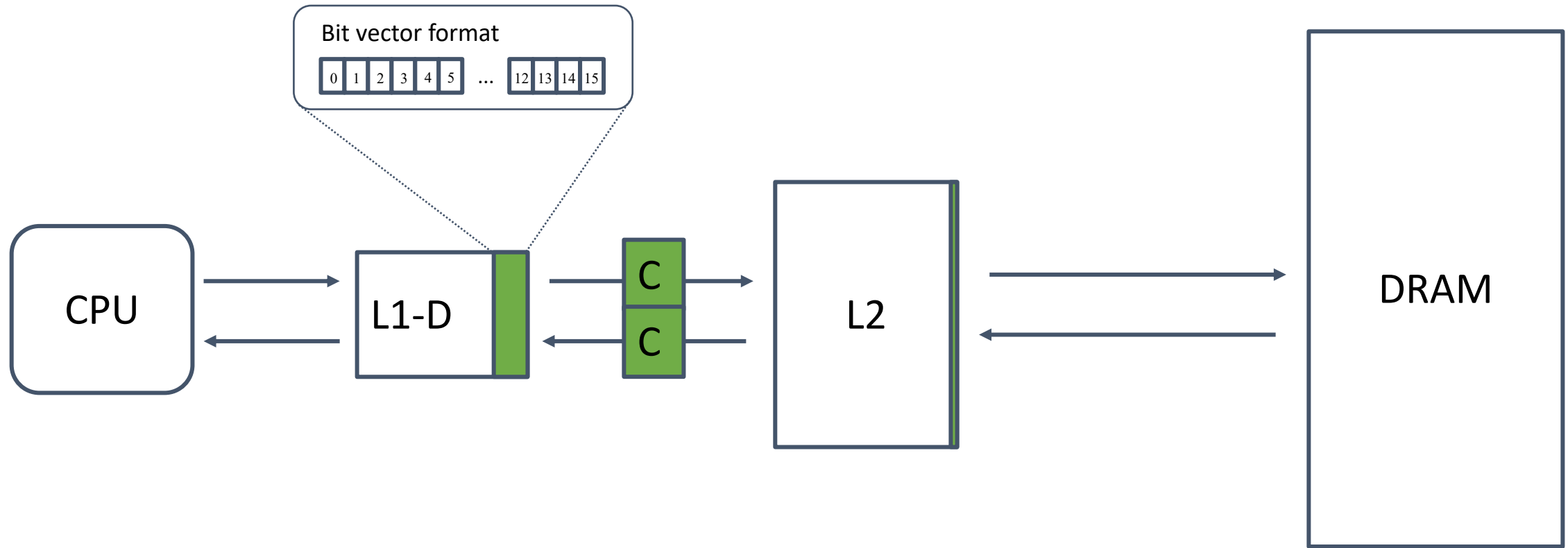


Microarchitectural Overview

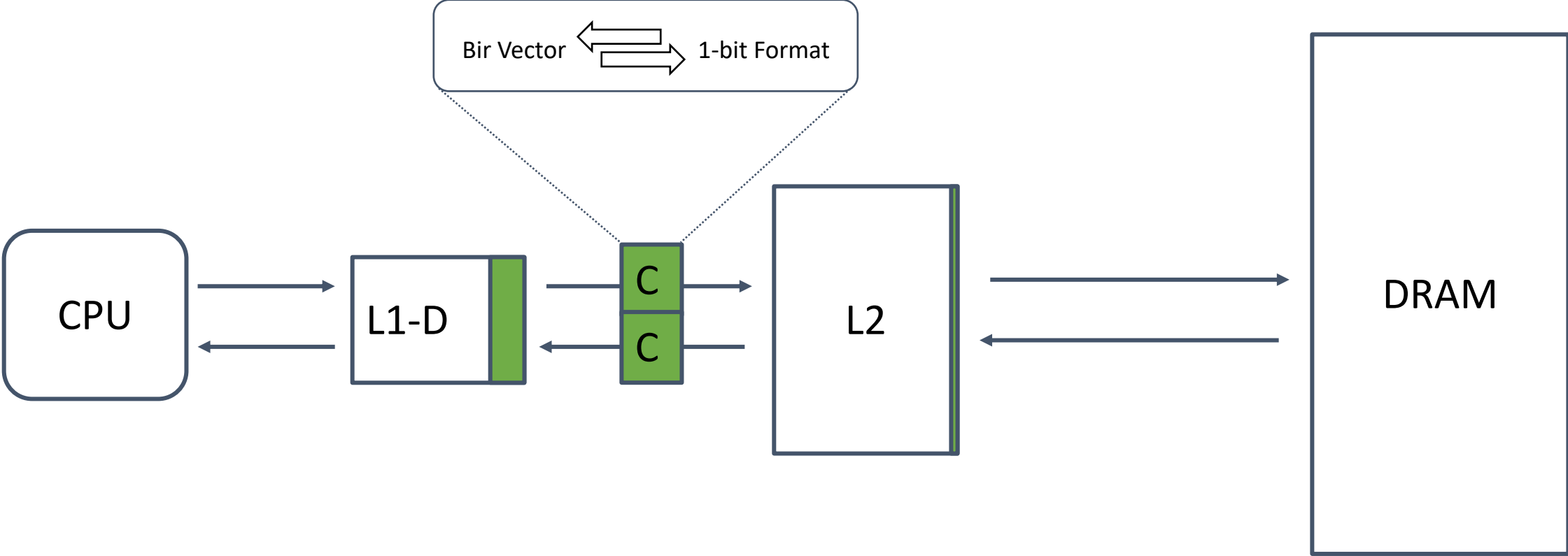
ZeRØ: Microarchitectural Overview



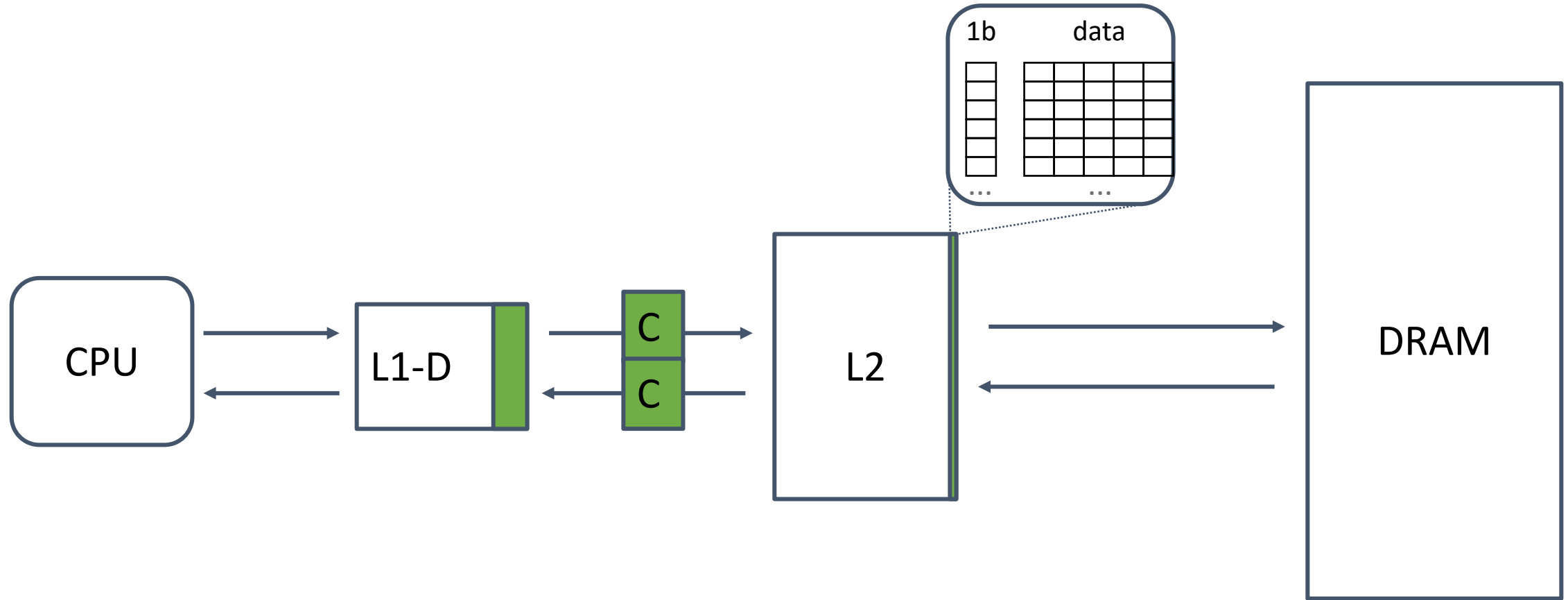
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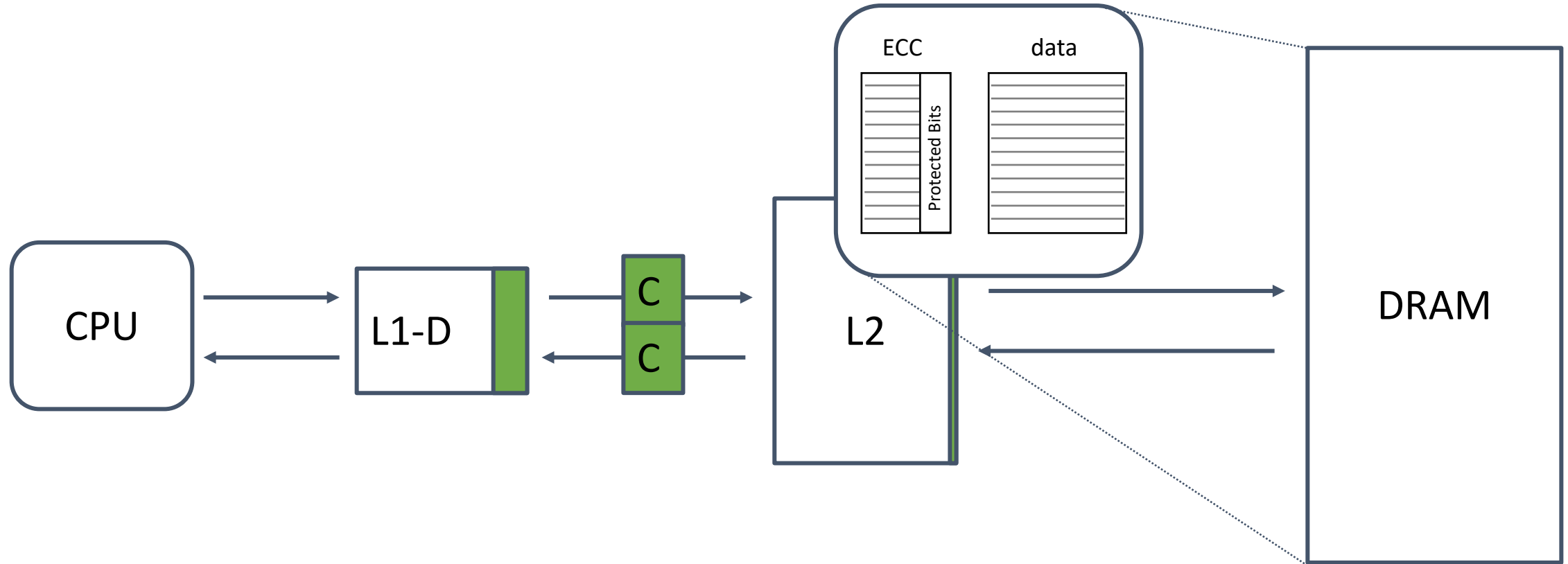
ZeRØ: Microarchitectural Overview



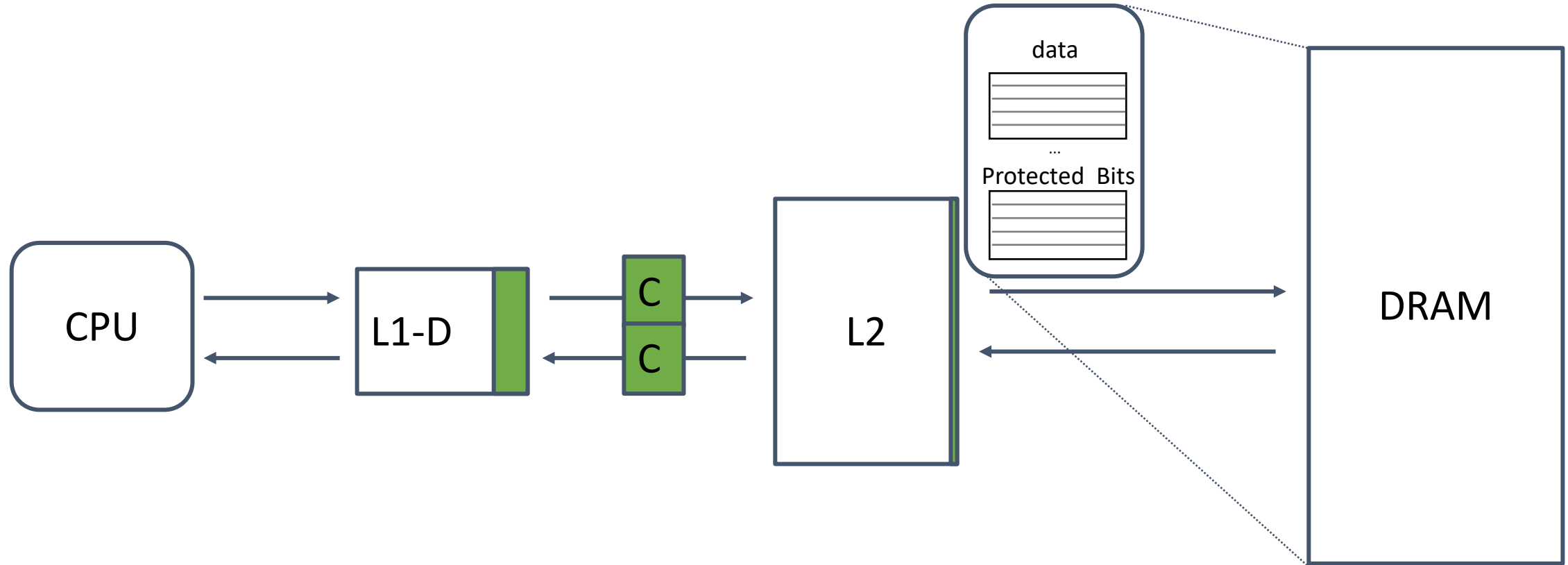
ZeRØ: Microarchitectural Overview



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ZeRØ: Microarchitectural Overview



ISA Extensions

ZeRØ: ISA Extensions

CPtrST/CPtrLD Address, Value

DPtrST/DPtrLD Address, Value

ZeRØ: ISA Extensions

CPtrST/CPtrLD Address, Value

DPtrST/DPtrLD Address, Value

*Same Layout as
regular Loads/Stores*

ZeRØ: ISA Extensions

CPtrST/CPtrLD Address, Value

DPtrST/DPtrLD Address, Value

ClearMeta Address, Mask

*Only invoked upon
free() or delete()*

Performance

ZeRØ: Performance Overheads

➤ Hardware Overheads.

➤ Software Overheads.

ZeRØ: Performance Overheads

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- Our hardware measurements show that ZeRØ has minimal latency/area/power overheads.

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ZeRØ: Performance Overheads

➤ Hardware Overheads.

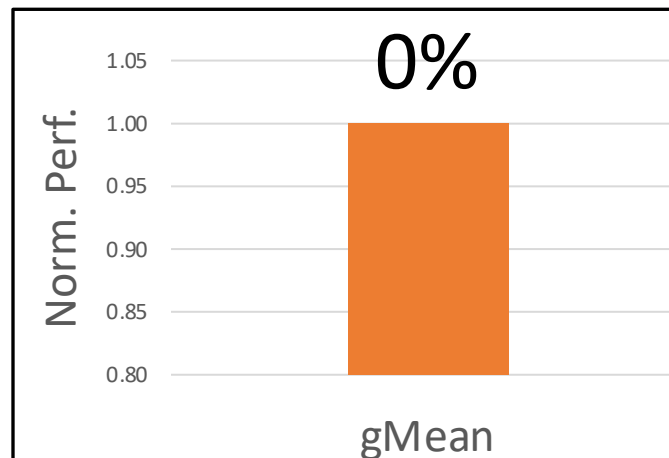
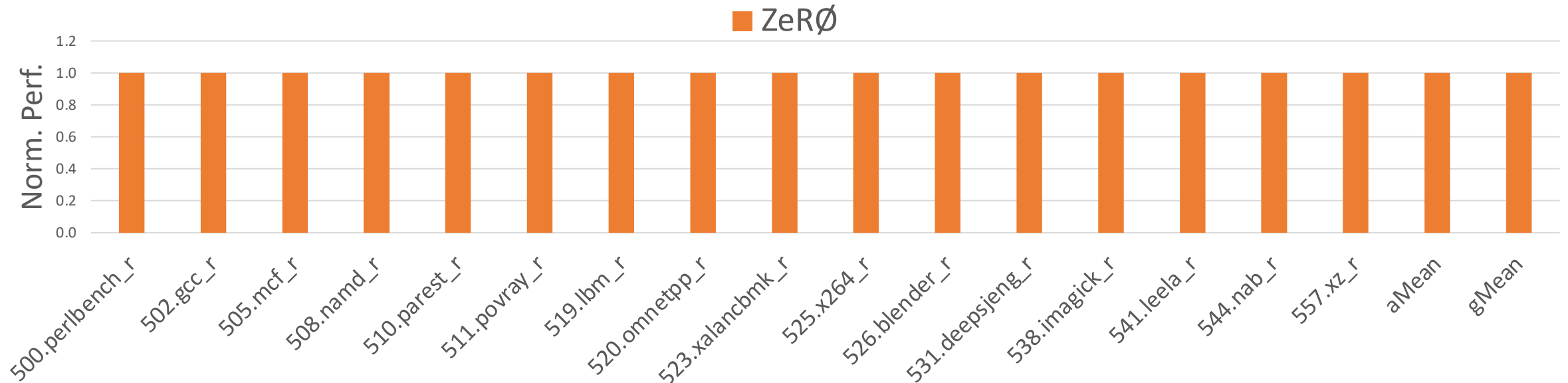
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➤ Software Overheads.

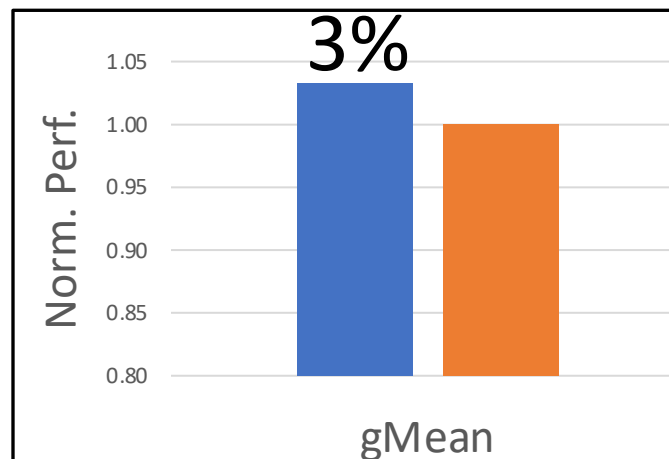
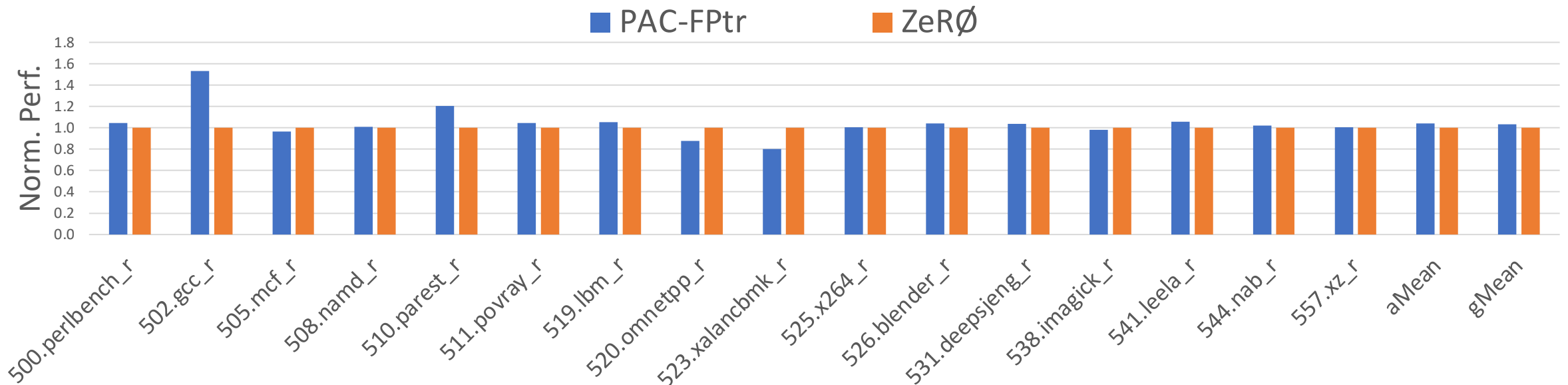
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The ClearMeta instructions are emulated on x86_64 using dummy stores

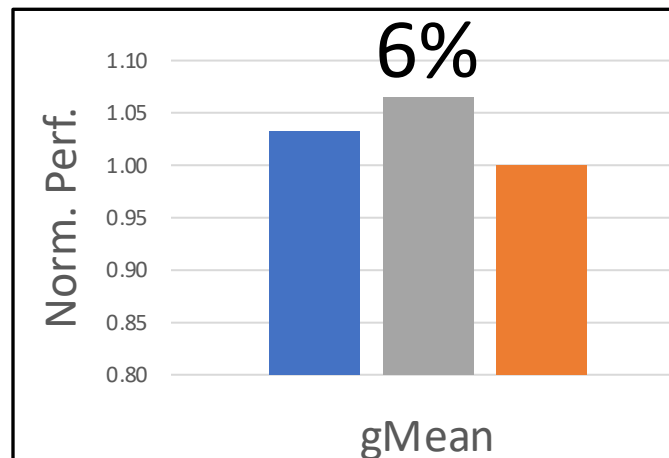
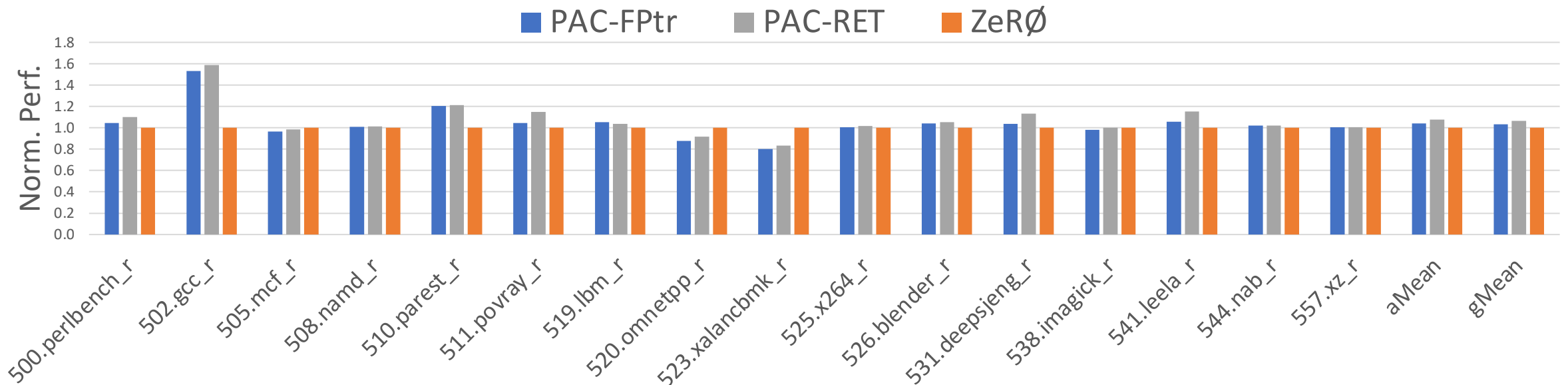
ZeRØ: Performance Results on x86_64



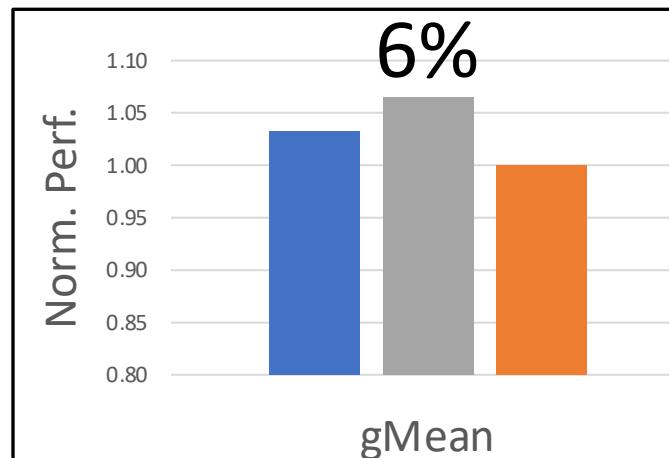
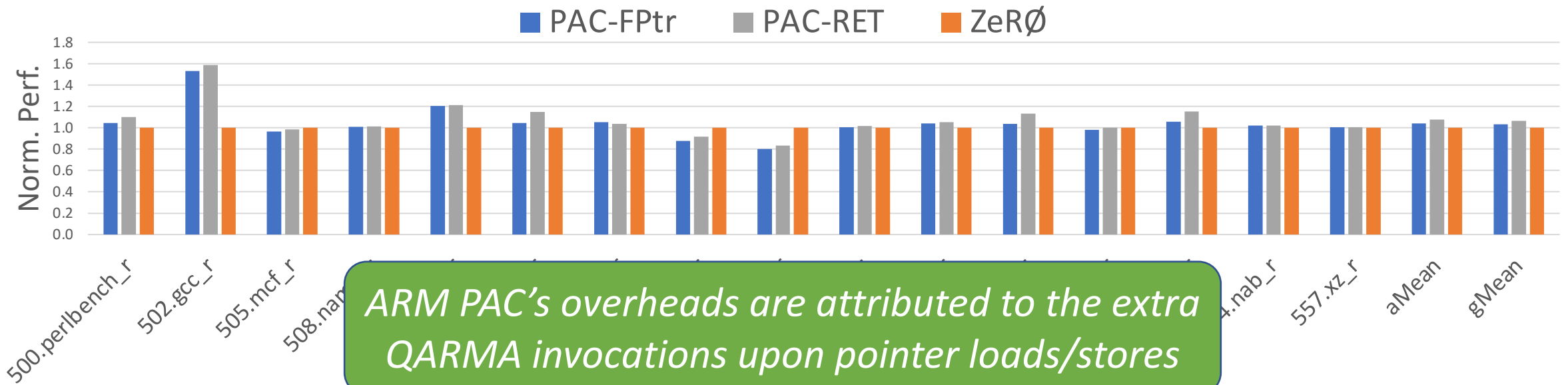
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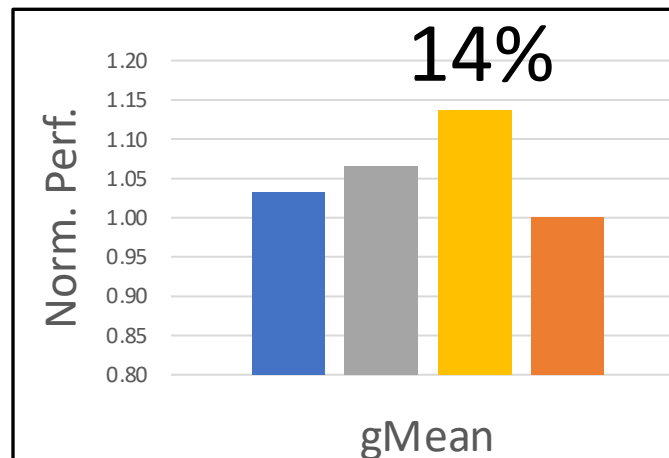
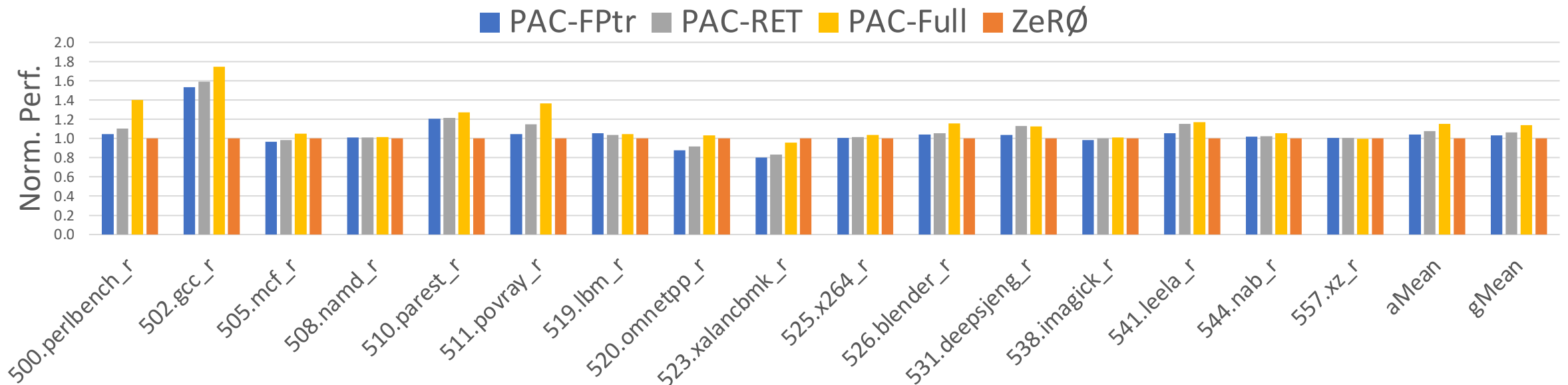
ZeRØ: Performance Results on x86_64



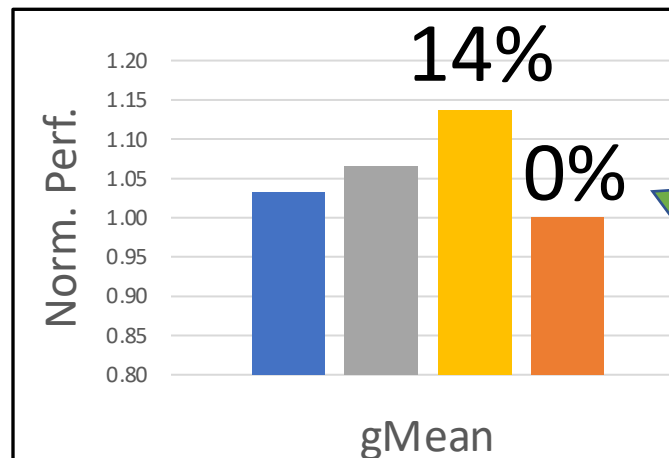
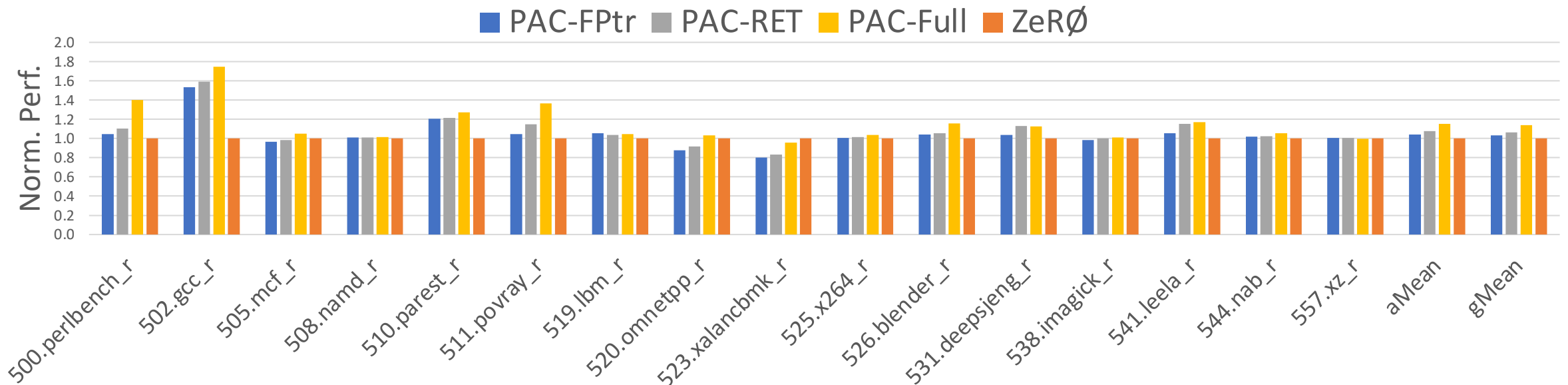
ZeRØ: Performance Results on x86_64



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ZeRØ: Performance Results on x86_64



ZeRØ reduces the average runtime overheads of pointer integrity from 14% to 0%

Limitations

ZeRØ: Limitations

- Non-pointer data corruption attacks.
 - Require a full memory safety solution.

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- Non-pointer data corruption attacks.
 - Require a full memory safety solution.
- Third-party code.
 - Clear the metadata bits before passing pointers to shared libraries.

Conclusion

- ZeRØ provides an efficient pointer integrity mechanism:
 - Is **easy to implement**.
 - Has **no runtime overheads**.
 - Offers **robust security**.
- ZeRØ can be applied to a wide variety of systems:
 - Ranging from servers to mobile devices.

The logo for ZeRØ, featuring the text "ZeRØ" in a white, sans-serif font. The letter "Ø" is stylized with a diagonal slash through it. The logo is positioned on the right side of the slide, set against a dark blue background that has a torn-paper edge effect on its left side.